

# Unit - I

## Security concepts:-

### Introduction:- and Need for security:-

- \* Information security is also called as cryptography and network security
- \* It is about how to secure our data from third party.
- \* whenever we are sending information to a friend (or) Receiver, we should make sure that the information is delivered safely
- \* we should make sure that the information is delivered to the Receiver without modifications.



\* here 'S' stands for sender and 'R' stands for Receiver

\* The communication b/w the sender and Receiver will obviously take place through internet.

\* Whenever we are sending information to Receiver, we should make sure that no third party will be having access to this information

\* If any third party is having access to the information that you are sending to the Receiver then the data is corrupted.

\* The corrupted data is nothing but, the data may be change (or) confidentiality of the data may be lost.

\* If you don't maintain the security, there is a chance that your data may be hacked.

\* For example:- If you and your friend wants to meet at 2:00pm. But you send a text message to your friend that to meet at 2:00pm.

\* If that data being read by third person and he modify the data, that to meet at 4:00pm.

\* Instead of 2:00pm, he made it 4:00pm, and it is delivery to the Receiver at 4:00pm.

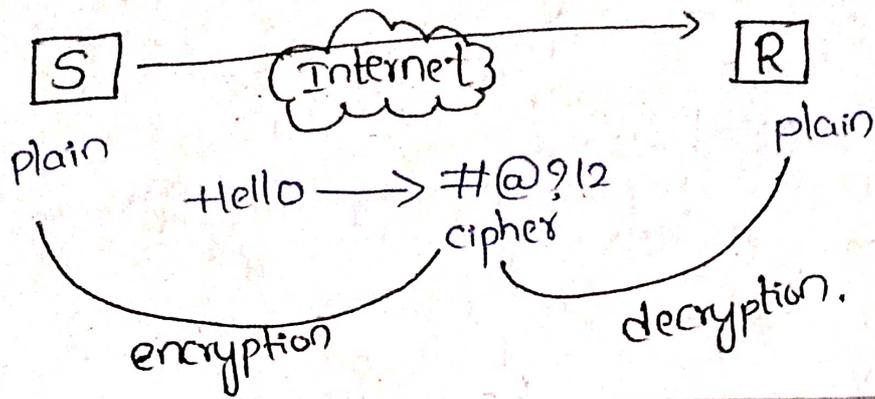
\* There miscommunication takes place and you both didn't meet.

\* Whenever the sender we are sending information from sender to the Receiver, two process will takes place i.e.,

- i) encryption
- ii) decryption.

\* Encryption:- It converting plain text (Hello) to cipher text (#@?12) (unreadable text).

\* Decryption:- It converts cipher text to plain text.



## Security Approaches:-

\* There are three ways that we can approach the security.

- 1) prevention
- 2) protection
- 3) Resilience

1) prevention:- It will prevent the threats by identifying the underlying causes before they occur.

\* It happens before the occurrence of threats.

2) protection:- It takes places, when the threats are ready to occur.

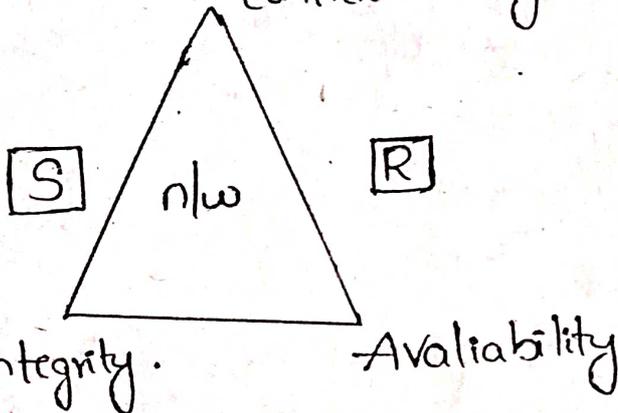
3) Resilience:- Here, the threat will already occur. When we are not in a position to control a threat then we have to adopt a mechanism (or) method (or) write a program through which the threat can be solved.

\* This is about security approaches.

#### 4) principles of security:-

- \* we need security to satisfy the confidentiality, integrity and availability.
- \* It is also called as CIA Triad (three things).
- \* whenever we are sending information from sender to Receiver, we have to maintain this CIA Triad for a proper and reliable communication.

\*



#### Confidentiality:-

- \* Confidentiality is nothing but confidential data (or) confidential message. should be kept in secret.
- \* The ~~data~~ information whenever we are sending from sender to Receiver, should be known only to the sender and Receiver not to any other third party.

#### Integrity:-

Integrity is nothing but whatever the data we are sending from sender to Receiver, it should send to the Receiver without any modifications.

- \* The data should be send from sender side to Receiver-side without any modifications.



(without any modifications)

### Availability:-

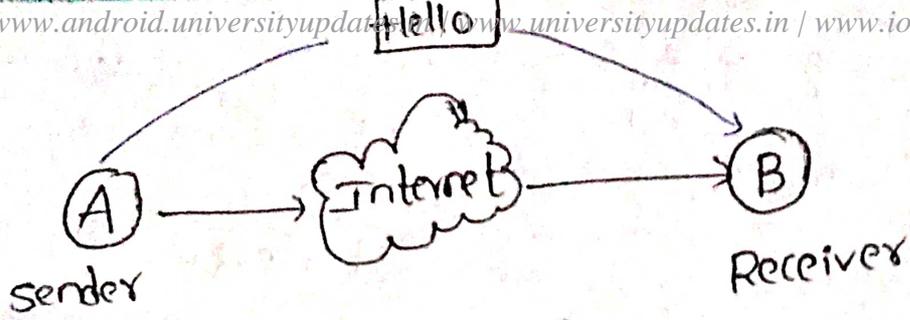
- \* Availability is nothing but whatever the data we are sending from sender to the Receiver, it should be available in all forms.
- \* The Receiver should be able to read the data and write the data, execute the data, modify the data.
- \* Receiver should be able to do each & every function.
- \* These are known as principles of security (or) goals of security. (or) maintaining the security to achieve CIA Triad!

### 5) Types of security Attacks:- (Possible ways of Attack)

- \* Any action that compromises the security of information.
- Types :- i) passive attack (read)  
ii) Active attack. (read, write, modify etc.)

### Passive attack:-

Whenever the data sending from the sender to Receiver, the third party can only read the data and observe the data without any modifications.



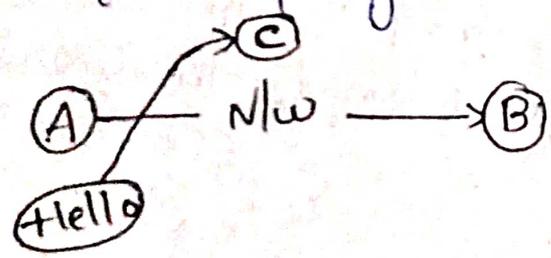
(C) (read).  
Third party

\* There are divided into two categories:-

- i) Release of message contents.
- ii) Traffic Analysis.

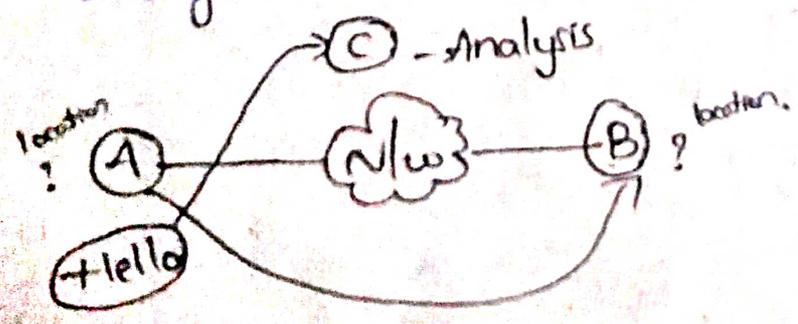
i) Release of message contents:- (Disclosure)

\* Whenever the data we are sending from sender to Receiver, the data will be released to third party also.



ii) Traffic Analysis:-

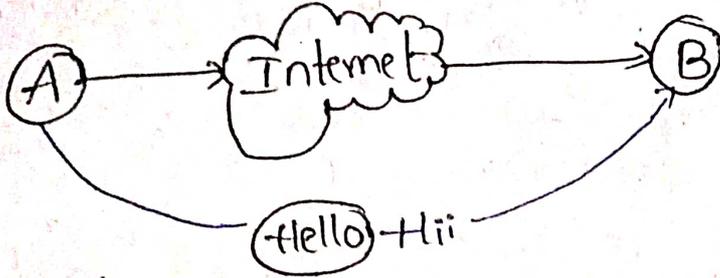
\* Whenever the data we are sending from sender to Receiver, third party try to observe and analyze the movement of the data.



# Active Attack:-

\* Whenever the data, we are sending from sender to Receiver, the third party can read, write, modify the data.

(C) Third party



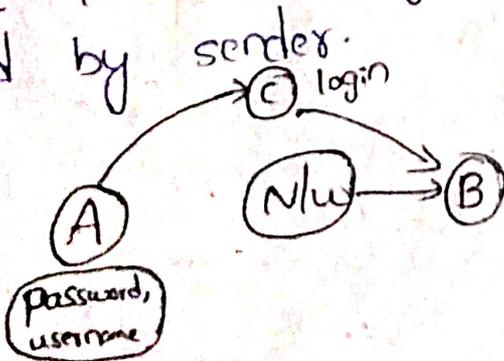
It can be divided into three categories:-

- i) Masquerade
- ii) Relay
- iii) Denial of service

## i) Masquerade:-

Whenever the data, we are sending from sender to receiver, the third party will stolen the data and it modify the data and sends to the Receiver.

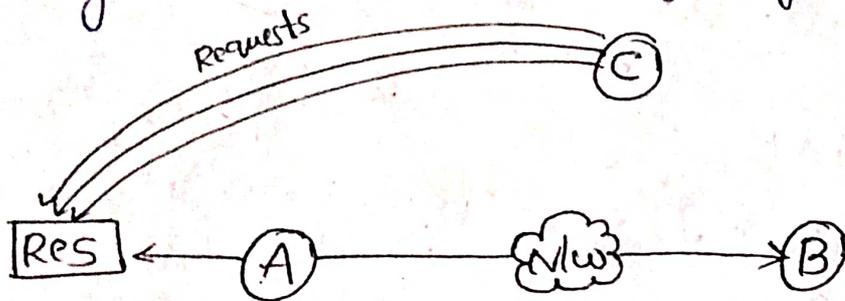
\* But Receiver thought that, the data is send by sender.



ii) Relay:- Whenever the data, we are sending from sender to Receiver, <sup>through the network</sup> the third party can read, write, modify the data.  
(Same as active attack).

ii) Denial of service:-  
 whenever sender wanted a Resource, third party wantedly send multiple Requests to the Resource.  
 \* In that case, Running capacity of Resource will be slow down.

\* Then sender has to wait and suffer.  
 \* Finally sender will be getting loss.



### 6) Security services:-

The services provided by security are:-

- 1) Authentication (user, password) (phone, otp)
- 2) Authorization (access control)
- 3) Non-Repudiation
- 4) Auditing. - Analyse



#### 1) Authentication:-

\* Getting an official permission to get into the website. and get into the server to access it.

\* There are many ways to check the authentication by <sup>checking</sup> whether (they are matching <sup>with</sup> their data) the username and password which you are using giving as an input is correct (or) not.

\* If the data is matched then you will be authenticated to use to the services.

## 2) Authorization:-

- \* After you are allowed to the enter into the website, upto what extent you can use this services of the server.
- \* It is also called as access control.
- \* It has some limitations that upto what extent you can use this services of the server.

## 3) Non Repudiation:-

- \* Once the message is transmitted from sender to Receiver
- \* Sender can't say that "No, I didn't send the message" as well as Receiver also
- \* This is also called as Non Repudiation.

Ex:- Money Transactions.

## 4) Auditing:-

- \* It will analyse the data, it will have entire information about the data
- \* If any unauthorization permissions happens then the Auditing will track the hacker.

## 7) security Mechanisms:-

To ensure the security we have some mechanisms

- i) Encipherment
- ii) Digital signature
- iii) Access control
- iv) Authentication Exchange
- v) Traffic padding
- vi) Routing control.

i) Encipherment:- (hide)

- \* The data will be hidden by cipher
- \* The sender will convert the data into a unreadable format means sender hides the data
- \* When the Receiver, Receives the data which is in unreadable that is converted into readable format.

ii) Digital signature:-

- \* Some special identity which is used for authentication.
- \* It is like a thumbnail and stamp
- \* It is also used for integrity of data.

iii) Access control:-

- \* Restricting the permissions to several levels.
  - \* In any organization, upto what extent of permissions can be given to a particular persons
- ex:- college management.

iv) Authentication exchanger:-

- \* Declaring the user as an authenticated user by comparing the username and password with the data that we are having in database.
- ex:- login instagram.

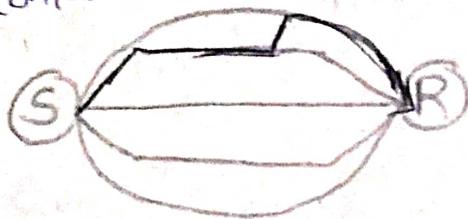
v) Traffic padding:-

- \* We have to add extra bits in the beginning (or) in the middle (or) in the ending in order to confuse the observer. (or) hacker.



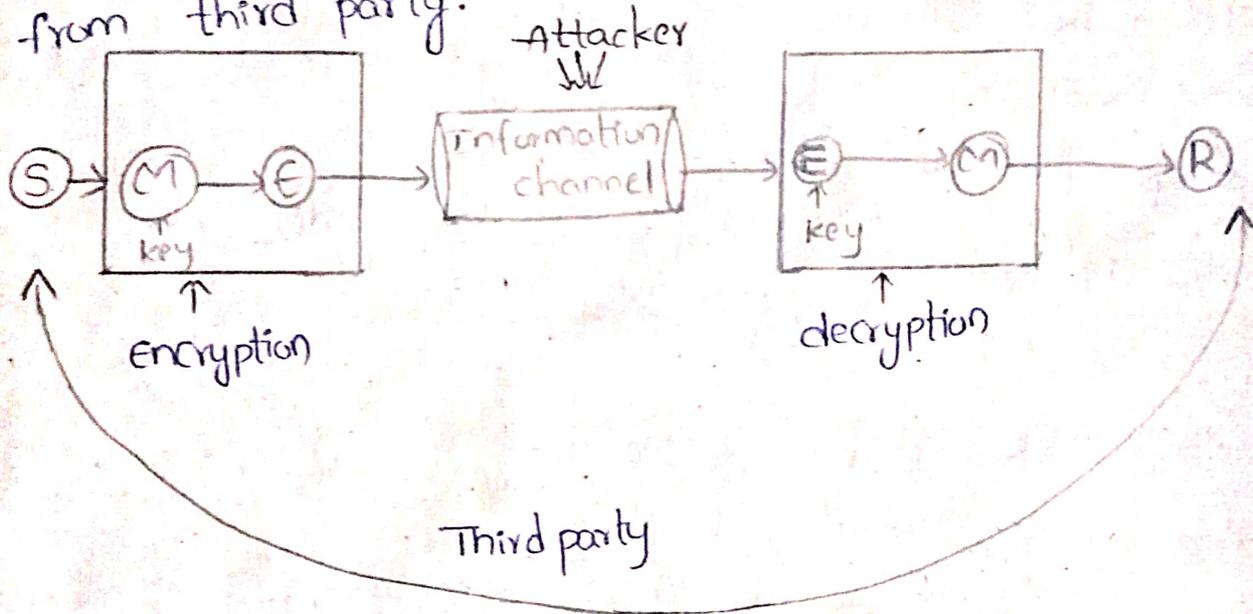
### vi) Routing control

- \* we have 'n' number of paths, we can go with any path, that is our wish.
- \* we can go with a mixture of path in order to confuse the hacker



### 8) Model for Network Security:-

- \* This is about, whenever the data is sending from sender to Receiver without any attacks from third party.



- \* The sender will generate a message, that message will be converted into a encrypted message. by using a key, this process is called encryption.

- \* After encryption, the encrypted message will be passed into information channel.
- \* The Information channel acts as the medium for both sender and receiver.

- www.android.universityupdates.in / www.universityupdates.in / www.ios.universityupdates.com
- \* Through this medium only the sender and Receiver will sharing the data
  - \* In this area, there are many attackers etc to hack the data. So, we should be careful in that area.
  - \* After crossing the information channel, the encrypted message will come out of the information channel.
  - \* The encrypted message is converted into original message by using a key, this process is called as decryption
  - \* The converted message will be read by Receiver
  - \* The file have a trusted third party which provides a key for encryption and decryption ~~party~~ process.
  - \* This is about network security model.

## Part-B

Plain text and cipher text:-

plain text:-

- \* It refers to anything which humans can understand.
- \* This is may be as simple as English sentences (Hello), a script (or) java code.
- \* If you can make sense of what is written then it is in plain text.

cipher text:  
\* cipher text (or) encrypted text, is a series of randomized letters (hanklmn) and numbers which humans cannot make any sense.

\* An encryption algorithm takes in a plaintext message, runs the algorithm on the plaintext and produces a ciphertext.

\* The ciphertext can be reversed through the process of decryption, to produce the original plaintext.

Ex:-

\* we will encrypt a sentence using caesar cipher

\* plaintext: This is srija

\* ciphertext: Aopz pz wshpuale.

## 2) Substitution Techniques and Transposition Techniques (classical encryption Techniques)

Substitution Techniques:-

\* Replacing the plain text alphabets (or) digits (or) symbols with some other alphabets (or) digits.  
\* This is also called as Replacement.

Ex:- FREE → XYZA

\* There are six techniques

- i) caesar cipher
- ii) Monoalphabetic
- iii) playfair cipher
- iv) hill cipher
- v) one time pad
- vi) polyalphabetic

ii) Caesar cipher :-

\* converting the plain text into cipher text by using formula.

$$C = E(3, P) = (P + 3) \text{ mod } 26$$

\* converting the cipher text into plain text by using formula

$$P = D(3, C) = (C - 3) \text{ mod } 26$$

*	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
PT →	A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z
CT →	d	e	f	g	h	i	j	k	l	m	n	o	p	q	r	s	t	u	v	w	x	y	z	a	b	c

Ex:- TROUBLE FREE [PT → CT]  
wuxxeh iuhh

Ex:- wuxxeh iuhh [CT → PT]  
TROUBLE FREE

ii) Monoalphabetic cipher :-

- \* Monoalphabetic means only one alphabet
- \* It has one-one relationship.
- \* there single ciphertext for each plaintext

Ex:- ALWAYS [PT → CT]  
VxA@k

Note:- we have to use only one alphabet for the same alphabet in plaintext.

\* The disadvantage is; the hacker can easily decode it.

ii) polyalphabetic cipher:-

- \* polyalphabetic means many alphabets.
- \* It has many-one relationship.
- \* there many ciphertext for a plaintext.

ex:- ALWAYS  
 (K) O Y (T) T P

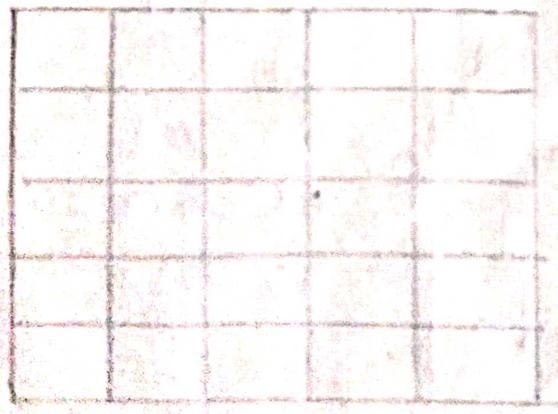
- \* we can use many other alphabets for the same alphabet in plaintext.

iii) playfair cipher:-

- \* It is also called as multiple letter encryption cipher
- \* there we have the plain text of msg & keyword we have to convert it to cipher text.
- \* we have some steps :-
  - 1) construct 5x5 matrix - 25 cells
  - 2) Fill the matrix
  - 3) Divide the msg → 2 letter pairs
  - 4) Apply rules & encrypt

Ex:- plain text = instruments ; key = monarchy.

step 1:-



Step 2:-

M	O	N	A	R
C	H	Y	b	d
e	f	g	i/j	k
l	p	q	s	t
u	v	w	x	z

(We have only 25 cells but we have 26 alphabets. So, we have to take 'i,j' as in one cell).

Step 3:- msg:- Instrument/sz

If any letter is remaining as odd then we have to add any alphabet.

Step 4:-

Rules:-

- i) If two alphabets are in same rows then Row → Right.
- ii) If two alphabets are in same columns then immediately goto column → down.
- iii) If two alphabets are not in same in row (or) column then draw a imaginary rectangle. Then we have to take corresponding horizontal alphabet.

M	O	N	A	R
C	H	Y	b	d
e	f	g	i/j	k
l	p	q	s	t
u	v	w	x	z

in → ga  
 st → tl  
 ru → mz  
 me → cd  
 nt → yq  
 sz → tx

n) Hill cipher:-

Encryption:-

To encrypt the message we have some steps:-

- i) construct the square matrix which is related to key matrix.
- ii) Assign the PT numbers to PT alphabets.

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25
A	B	C	D	E	F	G	H	I	J	K	L	M	N	O	P	Q	R	S	T	U	V	W	X	Y	Z

iii) By using the rule we are encrypt the msg

i.e.,  $C = Kp \pmod{26}$

Ex:-

key = VIEW ; message = ATTACK.

key matrix =  $\begin{bmatrix} V & E \\ I & W \end{bmatrix}$  plaintext matrices =  $\begin{bmatrix} A \\ T \end{bmatrix}, \begin{bmatrix} T \\ A \end{bmatrix}, \begin{bmatrix} C \\ K \end{bmatrix}$

$$\begin{bmatrix} 21 & 4 \\ 8 & 22 \end{bmatrix}$$

$$\begin{bmatrix} 0 \\ 19 \end{bmatrix}, \begin{bmatrix} 19 \\ 0 \end{bmatrix}, \begin{bmatrix} 2 \\ 10 \end{bmatrix}$$

→ Take the key as CDDF for easy calculation.

key =  $\begin{bmatrix} 2 & 3 \\ 3 & 6 \end{bmatrix}$

now apply  $Kp \pmod{26} = C$

①  $\begin{bmatrix} A \\ T \end{bmatrix} = \begin{bmatrix} 2 & 3 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} 0 \\ 19 \end{bmatrix} \pmod{26}$

$$= \begin{bmatrix} 2(0) + 3(19) \\ 3(0) + 6(19) \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 57 \\ 114 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 5 \\ 7 \end{bmatrix}$$

$$= \begin{bmatrix} F \\ K \end{bmatrix}$$

**AT → FK**

$$2) \begin{bmatrix} T \\ A \end{bmatrix} = \begin{bmatrix} 2 & 3 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} 19 \\ 0 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 2(19) + 3(0) \\ 3(19) + 6(0) \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 38 \\ 57 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 12 \\ 5 \end{bmatrix}$$

$$= \begin{bmatrix} M \\ F \end{bmatrix}$$

**TA → MF**

$$3) \begin{bmatrix} C \\ K \end{bmatrix} = \begin{bmatrix} 2 & 3 \\ 3 & 6 \end{bmatrix} \begin{bmatrix} 2 \\ 10 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 2(2) + 3(10) \\ 3(2) + 6(10) \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 34 \\ 66 \end{bmatrix} \pmod{26}$$

$$= \begin{bmatrix} 8 \\ 14 \end{bmatrix}$$

$$= \begin{bmatrix} I \\ O \end{bmatrix}$$

**CK → IO**

**ATTACK  $\rightarrow$  FKMFIO**

Decryption:-

\* By using the rule we can decrypt the msg.

$$P = K^{-1} C \text{ mod } 26$$

$$K^{-1} = \frac{1}{|K|} \text{adj } K$$

Eg:- cipher text = (ATTACK)FKMFIO

plain text = ATTACK

\* Determinant of matrix  $D = \begin{vmatrix} a & b \\ c & d \end{vmatrix}$

$$D = |ad - bc|$$

Eg:-  $K = \begin{vmatrix} 2 & 3 \\ 3 & 6 \end{vmatrix} = |12 - 9| = 3$

$\therefore$  determinant value of  $K = 3$

\* Multiplicative inverse of determinant is;

•  $dd^{-1} = 1 \text{ mod } 26$

•  $dd^{-1} \text{ mod } 26 = 1$

Eg:-  $3 \times K^{-1} = 1 \text{ mod } 26$

$3 \times K^{-1} \text{ mod } 26 = 1$

$3 \times 9 \text{ mod } 26 = 1$

$27 \text{ mod } 26 = 1$

so,  $K^{-1} = 9$

\* Adjacent matrix of A

Let  $A = \begin{bmatrix} a & b \\ c & d \end{bmatrix}$  then  $\text{Adj}(A) = \begin{bmatrix} d & -b \\ -c & a \end{bmatrix}$

eg:-  $K = \begin{bmatrix} 2 & 3 \\ 3 & 6 \end{bmatrix}$ ,  $\text{adj}(K) = \begin{bmatrix} 6 & -3 \\ -3 & 2 \end{bmatrix}$

\* Before decryption, we have to remove negative values. By adding 26 to negative values.

$\therefore \text{adj}(K) = \begin{bmatrix} 6 & -3+26 \\ -3+26 & 2 \end{bmatrix} = \begin{bmatrix} 6 & 23 \\ 23 & 2 \end{bmatrix}$

Now,  $K^{-1} = \frac{1}{|K|} \text{adj}(K)$

$\therefore \frac{1}{|K|} = |K^{-1}|$

$K^{-1} = 9 \begin{bmatrix} 6 & 23 \\ 23 & 2 \end{bmatrix}$

$K^{-1} = \begin{bmatrix} 54 & 207 \\ 207 & 18 \end{bmatrix}$

\* We have to do mod 26 for simple calculation

$K^{-1} = \begin{bmatrix} 54 & 207 \\ 207 & 18 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 2 & 25 \\ 25 & 18 \end{bmatrix}$

Now, we will decrypt

cipher = FK MF IO

$C = \begin{bmatrix} F \\ K \end{bmatrix} = \begin{bmatrix} 5 \\ 10 \end{bmatrix}$

$\therefore \text{plain text } P = K^{-1} C \text{ mod } 26$

$P = \begin{bmatrix} 2 & 25 \\ 25 & 18 \end{bmatrix} \begin{bmatrix} 5 \\ 10 \end{bmatrix} \text{ mod } 26$

$P = \begin{bmatrix} 200 \\ 305 \end{bmatrix} \text{ mod } 26 = \begin{bmatrix} 0 \\ 19 \end{bmatrix} = \begin{bmatrix} A \\ T \end{bmatrix}$

Similarly,  

$$C = \begin{bmatrix} M \\ F \end{bmatrix} = \begin{bmatrix} 12 \\ 5 \end{bmatrix}$$

So, corresponding plain text is;

$$P = \begin{bmatrix} 2 & 25 \\ 25 & 18 \end{bmatrix} \begin{bmatrix} 12 \\ 5 \end{bmatrix} \pmod{26} = \begin{bmatrix} 149 \\ 396 \end{bmatrix} \pmod{26} = \begin{bmatrix} 19 \\ 0 \end{bmatrix} = \begin{bmatrix} T \\ A \end{bmatrix}$$

Again,

$$C = \begin{bmatrix} I \\ 0 \end{bmatrix} = \begin{bmatrix} 8 \\ 14 \end{bmatrix}$$

$$P = \begin{bmatrix} 2 & 25 \\ 25 & 18 \end{bmatrix} \begin{bmatrix} 12 \\ 5 \end{bmatrix} \pmod{26} = \begin{bmatrix} 366 \\ 452 \end{bmatrix} \pmod{26} = \begin{bmatrix} 2 \\ 10 \end{bmatrix} = \begin{bmatrix} C \\ K \end{bmatrix}$$

Finally, the plain text is ATTACK.

v) one-time pad :- / verman cipher

\* The condition is here;

The length of key should be equal to the length of plain text.

\* key length = length of PT.

Ex:- PT = security

key = Acmtkyiv

PT → security → 18 4 2 20 17 8 19 24  
 key → Acmtkyiv → 0 2 19 19 10 24 8 21

18	6	14	39	27	32	27	45	add
			-26	-26	-26	-26	-26	
18	6	14	13	1	6	1	19	sub
S	G	O	N	B	G	B	T	

\* SGIONBGIB is cipher text for security

\* This is encryption.

\* For decryption, do the same process in reverse.

### Transposition Techniques:-

\* This is also called as Rearrangement.

\* Rearrange the plain text alphabets (or) digits (or) symbols with the same plain text alphabet (or) digits (or) symbols which are given.

\* We shouldn't add any other alphabets.

\* Ex:- FREE → EREF

REEF

FEER etc...

\* There are four techniques

i) Rail Fence Transposition

ii) Columnar Transposition

iii) Improved Transposition

iv) Book cipher

### i) Rail Fence Transposition:-

\* We can Rearrange the plain text into cipher text by using the depth which is equal is 2.

Ex:- TROUBLE FREE

T O U B L E F R E  
R U L F E

CT → TOBERERULFE

PT → diagonal

CT → Row.

www.android.universityupdates.in / www.universityupdates.in / www.ios.universityupdates.in

- \* It is useful for short messages.
- \* It is not so efficient.

ii) Columnar Transposition: -

- \* We have to arrange the plain text into a matrix.
- \* It is not a mandatory to take a square matrix only.
- \* We can take any matrix like rectangle, square, etc...
- \* Fill the matrix with the plain text in a row wise.
- \* Eg: - Information security  $\rightarrow$  plain text

I	N	F	O	R
M	A	T	I	O
N	S	E	C	U
R	I	T	Y	

- \* Generate the key, which is in the form of number & which is less than <sup>(or) equal to</sup> the no. of columns we took.
  - \* Then write the corresponding <sup>cipher text</sup> column wise
- key = 32514.
- \* We have to select the key randomly

	1	2	3	4	5
I	n	f	o	r	
m	a	t	i	o	
n	s	e	c	u	
r	i	t	y		
↓	↓	↓	↓		

key = 32514

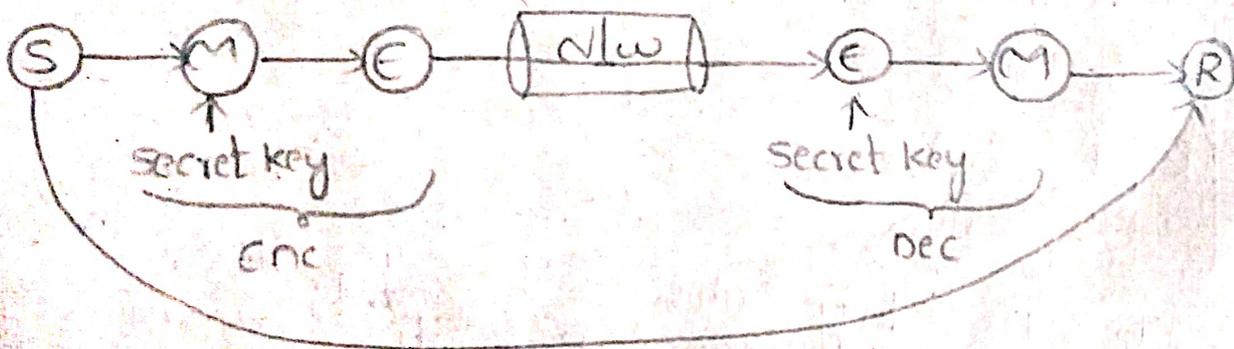
ftetn asivouimnroicy → information security.

#### 4) Symmetric and Asymmetric key cryptograph

Symmetric key cryptograph:-

\* We have only one key on sender side and Receiver side.

\* We are using only one key for encryption and decryption process.



\* sender wants to send a message to Receiver

\* sender generates the message after that, the message has to be encrypted with the help of secret key. This process is known as encryption.

\* The encrypted message will be enter into the network

\* After that the encrypted message will come out of the network.

\* Then, encrypted message is converted into original message with the help of same secret key which is used at the encryption process.

- \* This process is known as decryption.
- \* The original message is read by the receiver.
- \* The disadvantage is; it can easily implement because we have only one secret key.
- \* It is not so efficient.
- \* It is not at all secure.

### Asymmetric key cryptography:-

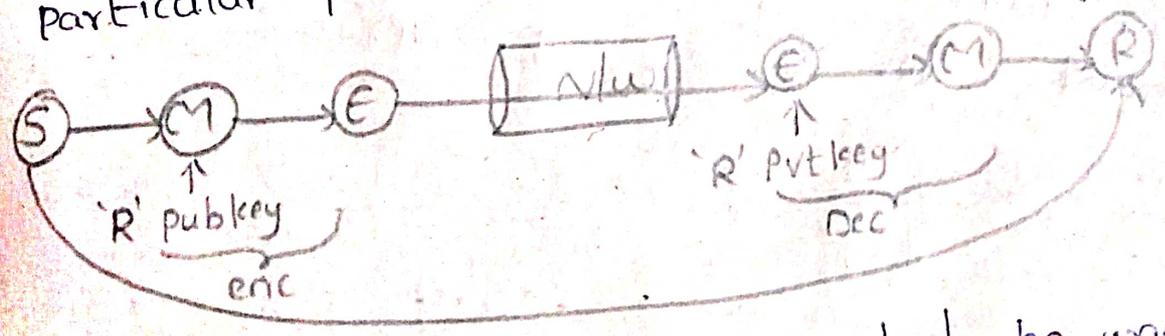
\* We have different keys on senderside and receiver side.

\* We have two types of keys:-

- 1) public key
- 2) private key.

\* public key is a key which is known to everyone.

\* private key is a key which is known to a particular person



\* (Sender generates the message, which he wants to send to a Receiver)

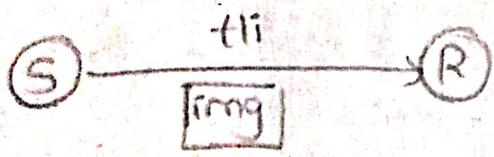
\* Sender wants to send a message to Receiver.

\* sender generates the message after that, the message has to be encrypted with the help of Receiver's public key. This process is

- \* The encrypted message will be enter into the network.
- \* The encrypted message will come out from the network.
- \* Then, encrypted message will be converted into original message with the help of Receiver's private key.
- \* This process is known as decryption.
- \* The original message will be read by Receiver.
- \* In this, we have more security when compar to symmetric key cryptography.

### 5] Steganography:-

- \* Hiding information with in another message.
- \* Embedding the msg with in an image, (or) video (or) pdf.
- \* After transferring the msg from sender to Receiver, later msg is extracted from embedded devices by Receiver.



- \* We have several steganography techniques:-
  - 1) Least significant bit (LSB)
  - 2) Audio/video steganography
  - 3) character marking etc...

\* we have some attacks in steganography. like the hacker will observe the data and modify the data.

6] key size and key range