

# GRAPH THEORY.

## \* Definition of a Graph

A Graph  $G$  is a finite non-empty set of vertices  $V$  with a set of edges  $E$  which consists of pairs of vertices.

ie.,  $G = (V, E)$  where

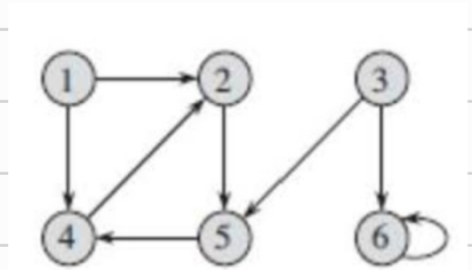
$V \rightarrow$  Set of vertices

$E \rightarrow$  set of edges of the form  $(x, y)$   
such that  $x, y \in V$ .

## \* Directed Graph:-

A graph in which edges have directions. Thus the edge set is a set of ordered pairs of vertices.

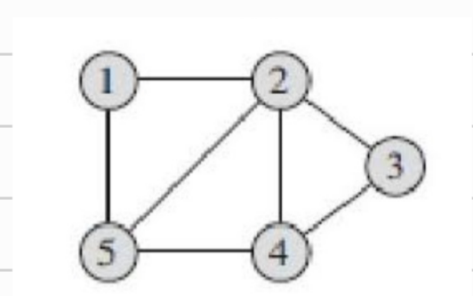
In this example, the vertex set  $V = \{1, 2, 3, 4, 5, 6\}$ ,  
edge set  $E = \{(1, 2), (1, 4), (2, 5), (3, 5), (3, 6), (4, 2), (5, 4), (6, 6)\}$



## \* Undirected Graph:-

A graph in which edges do not have directions.

In this example, vertex set  $V = \{1, 2, 3, 4, 5\}$ ,  
edge set  $E = \{(1, 2), (1, 5), (2, 3), (2, 4), (2, 5), (3, 4), (4, 5)\}$



Note i- Since directions are not given we can write an edge as  $(1, 2)$  or  $(2, 1)$ .  
ii) Vertices are also called as Nodes.

## \* Adjacent vertices

The vertices with which a given vertex forms an edge.

## \* Degree of a vertex

→ In undirected graph, degree of a vertex is the number of edges formed by it

→ In directed graph, if we have an edge from A to B, then A is called as the Source and B is called as Sink.

Here, indegree of a vertex is the no. of edges coming towards it (sink) and outdegree is the no. of edges going away from it. (source vertex)

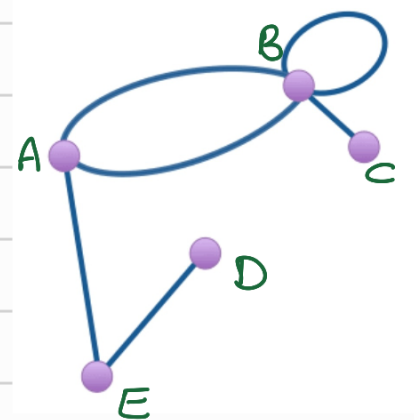
## \* Parallel Edges :-

multiple edges between same pair of vertices

## \* Loops :-

an edge from a vertex to itself.

Ex :- Here, we have parallel edges between A and B and there is a loop at vertex B.

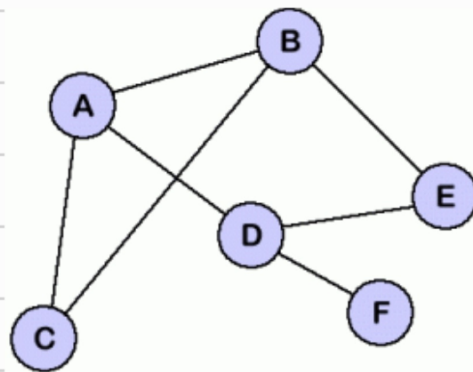


Degrees of the vertices:-

$$\text{deg}(A) = 3, \quad \text{deg}(D) = 3$$

$$\text{deg}(B) = 3, \quad \text{deg}(E) = 2$$

$$\text{deg}(C) = 2, \quad \text{deg}(F) = 1$$



Indegree  $\rightarrow d^+$ , outdegree  $\rightarrow d^-$

$$d^+(1) = 0$$

$$d^-(1) = 1$$

$$d^+(2) = 1$$

$$d^-(2) = 3$$

$$d^+(3) = 2$$

$$d^-(3) = 0$$

$$d^+(4) = 1$$

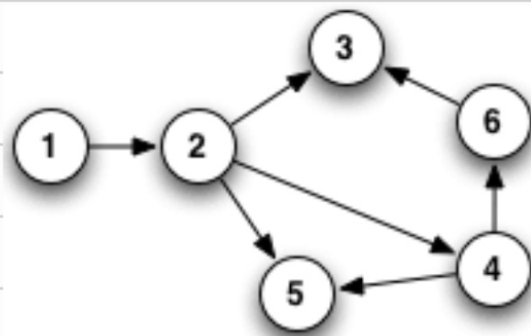
$$d^-(4) = 2$$

$$d^+(5) = 2$$

$$d^-(5) = 0$$

$$d^+(6) = 1$$

$$d^-(6) = 1$$



Theorem 1:- Handshaking Theorem

In a graph, the sum of degrees of all vertices is equal to twice the no. of edges.

Note i-) Sum of degrees of all vertices is always even.

ii) The no. of odd vertices in a graph is even

Types of Vertices:-

1. Isolated Vertex:- vertex with degree 0.

2. Pendant Vertex:- vertex with degree 1

3. Odd Vertex:- vertex with odd degree

4. Even Vertex :- vertex with even degree

5. Source/Sink vertex:- In a directed edge, A is the source vertex and B is the sink vertex.

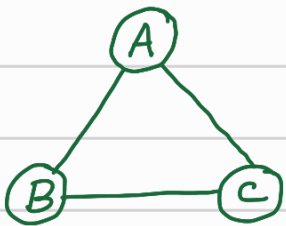


### Types of Graphs :-

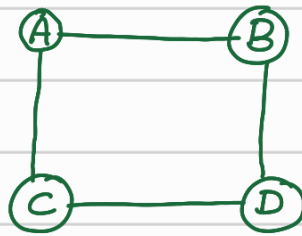
1. Simple Graph: A graph without parallel edges and loops.

2. Point Graph: A graph with only one vertex & 0 edges

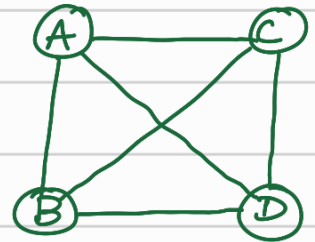
3. Regular Graph: A graph in which degree of all vertices are equal. to  $r$ . It is called  $r$ -regular graph.



2-regular graph



2-regular graph.



3-regular graph.

4. Connected Graph: A graph in which we have a path from one vertex to any other vertex.

5. Complete Graph: A graph in which each vertex is connected with every other vertex. ( $K_n$ )

It is a graph with maximum no. of edges. It is represented using  $K_n$ .

Note:- In a simple graph with  $n$  vertices, the maximum no. of edges is  $\frac{n(n-1)}{2}$

6. Planar Graph: A graph which can be drawn in a plane such that there are no intersecting edges.

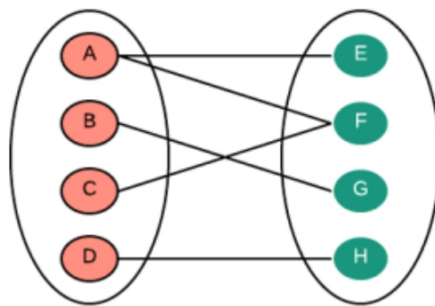
7. Bipartite Graph: A graph in which vertex set  $V$  is divided into two sub-vertex sets  $V_1$  &  $V_2$  and the edges are formed with one vertex from each  $V_1$  &  $V_2$ .

In this example,

$$V = \{A, B, C, D, E, F, G, H\}$$

$$V_1 = \{A, B, C, D\} \text{ \& } V_2 = \{E, F, G, H\}$$

edges are formed one vertex each from sets  $V_1$  &  $V_2$ .

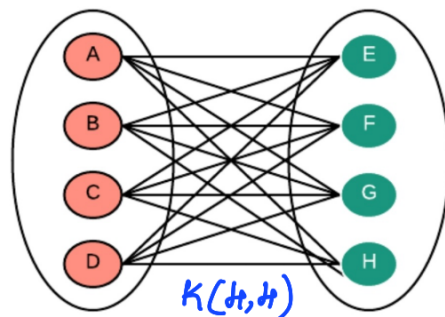


8. Complete Bipartite Graph: A bipartite graph in which each vertex from one sub-vertex set is connected with every vertex from second sub-vertex set.

It is represented as  $K(p, q)$  where

$p \rightarrow$  no. of vertices in  $V_1$

$q \rightarrow$  no. of vertices in  $V_2$ .



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- Compiled By Usha.